

Sequential Change-Point Approach for Online Community Detection

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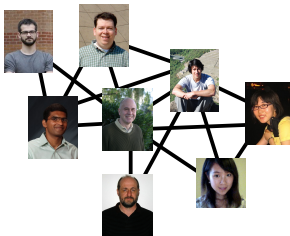


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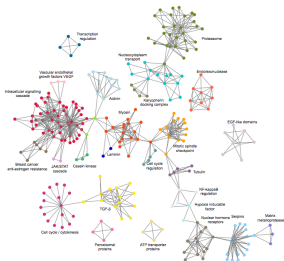
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Community

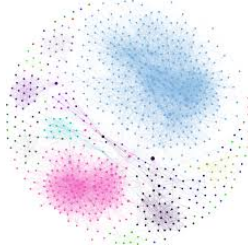
Collaboration network



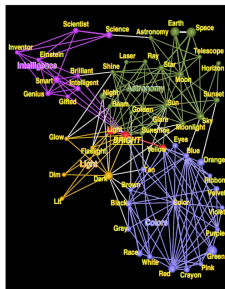
Protein interaction network



Facebook



Word association



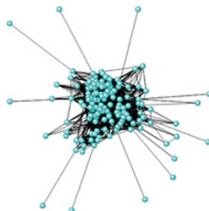
Enron email data set

Enron Emails Reveal What a Web of Deceit Really Looks Like

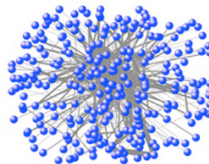
ALEXIS C. MADRIGAL | JUL 13 2011, 11:54 AM ET



The shape that a social network takes may be a new kind of digital smoke to spot the fires of corruption within an organization



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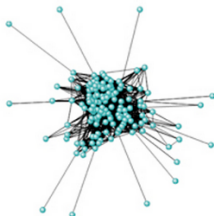
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Some networks might be structurally suspicious, even if none of the content passing on it looks that way ... diagnose bad acting within a large organization. – *The Atlantic*, 2011

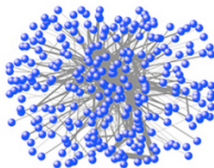
- ▶ 500,000 emails involving 151 unknown employees and more than 75,000 distinct addresses; each email with time stamp, sender and receiver
- ▶ between the years 1998 and 2002, record for 1,177 days

“legal project”: many people are connected to many others on a project, and information is widely distributed

“illicit”: information concentrated in a few hands



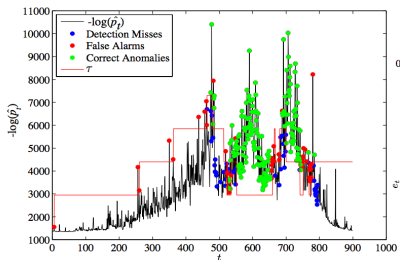
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Emergence of a community

Starting from a certain time, anomalous email discussion topics arise between a small group of people.



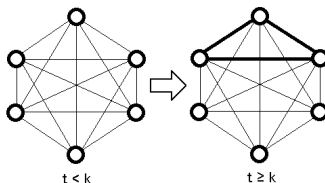
Date	Significance
Dec. 1, 2000	Days before “California faces unprecedented energy alert” (Dec. 7) and energy commodity trading deregulated in Congress. (Dec. 15) [37].
May 9, 2001	“California Utility Says Prices of Gas Were Inflated” by Enron collaborator El Paso [38], blackouts affect upwards of 167,000 Enron customers [39].
Oct. 18, 2001	Enron reports \$618M third quarter loss, followed by later major correction [40].

Online detection of community emergence

- ▶ a network with N nodes
- ▶ observe a **sequence** of independent adjacency matrices

$$X_1, X_2, \dots$$

- ▶ $X_i \in \mathbb{R}^{N \times N}$: interaction of nodes at time i
- ▶ there may exist an **unknown** time s.t. after that an **unknown** subset of nodes interact with higher frequency



- ▶ offline version [Arias-Castro-Verzelen2014]

Sequential change-point detection approach

- H_0 : X_i : Erdős-Renyi random graph

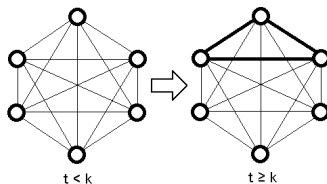
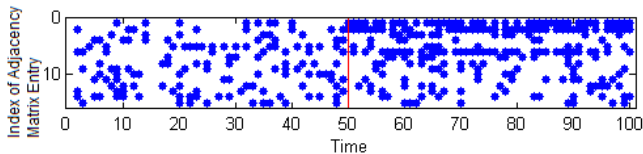
$$[X_t]_{ij} = \begin{cases} 1 & \text{w. p. } p_0 \\ 0 & \text{otherwise} \end{cases} \quad \forall (i, j).$$

- H_1 : there exists an unknown time κ such that afterwards **unknown** subset of nodes \mathcal{S}^* interact more frequently

$$[X_t]_{ij} = \begin{cases} 1 & \text{w. p. } p_1 \\ 0 & \text{otherwise} \end{cases} \quad \forall i, j \in \mathcal{S}^*, \quad t > \kappa,$$

$$[X_t]_{ij} = \begin{cases} 1 & \text{w. p. } p_0 \\ 0 & \text{otherwise} \end{cases} \quad \forall i \notin \mathcal{S}^* \text{ or } j \notin \mathcal{S}^*, \quad t > \kappa.$$

$$p_0 < p_1$$



- ▶ **Goal:** detect **emergence** of an unknown community as **quickly as possible**
- ▶ define a **stopping rule** T for sequential data such that
 - ▶ rarely raise false alarm when there is no change
 - ▶ raise alarm quickly after the change (small detection delay)

Classic change-point detection

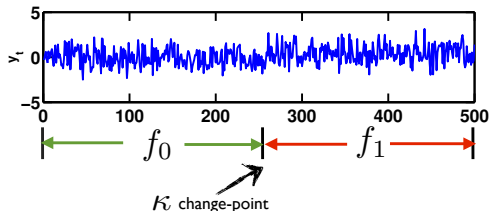
In statistics and quality-control

- ▶ Min-max formulation: Page (54), Lorden (71)
- ▶ Bayesian: Shirayev (63), Roberts (66)
- ▶ a sequence i.i.d. observations $y_1, y_2, \dots \in \mathbb{R}$
- ▶ unknown change-point $\kappa > 0$.

$$H_0 : y_t \sim f_0, \quad t = 1, 2, \dots$$

$$H_1 : y_t \sim f_0, \quad t = 1, \dots, \kappa,$$
$$y_t \sim f_1, \quad t = \kappa + 1, \dots$$

unknown change-point $\kappa > 0$



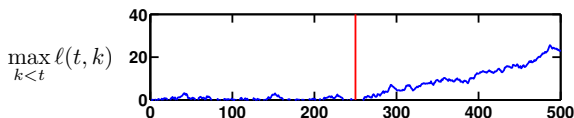
Likelihood ratio based procedure

- ▶ for a hypothesized $\kappa = k$:

$$\ell(t, k) = \log \frac{\prod_{i=1}^k f_0(y_i) \cdot \prod_{i=k+1}^t f_1(y_i)}{\prod_{i=1}^t f_0(y_i)} = \sum_{i=k+1}^t \log \frac{f_1(y_i)}{f_0(y_i)}$$

- ▶ likelihood ratio based change-point detection:

$$T = \inf\{t \geq 1 : \max_{k < t} \ell(t, k) \geq b\}$$



Normal distributions

► $f_0 = \mathcal{N}(0, 1)$, $f_1 = \mathcal{N}(\mu, 1)$, $\mu > 0$

► CUSUM procedure

$$T = \inf\{t : \max_{k < t} \sum_{i=k+1}^t (\mu y_i - \frac{\mu^2}{2}) \geq b\}$$

► when μ is **unknown**: $\hat{\mu}(k) = (\sum_{i=k+1}^t y_i)/(t - k)$
GLR procedure

$$T = \inf\{t : \max_{k < t} \frac{(\sum_{i=k+1}^t y_i)^2}{t - k} \geq b\}$$

Likelihood ratio based statistic

- ▶ for edge (i, j) , assumed change-point location $\kappa = k$, observation up to time t , likelihood ratio statistic given by

$$\begin{aligned} \ell(\kappa = k | \textcolor{red}{p}_1, \mathcal{S}) \\ = \sum_{(i,j) \in \mathcal{S}} \underbrace{\sum_{m=k+1}^t [X_m]_{ij} \log \left(\frac{p_1}{p_0} \right) + (1 - [X_m]_{ij}) \log \left(\frac{1 - p_1}{1 - p_0} \right)}_{U_{k,t,\textcolor{red}{p}_1}^{(i,j)}} \end{aligned}$$

- ▶ typically, we can assume p_0 **known** since it can be estimated from historic data
- ▶ $\textcolor{red}{p}_1$ is usually **unknown** since it represents anomaly

Exhaustive Search (ES) method

- ▶ Approach 1: assume unknown $p_1 = \delta$
- ▶ δ : nominal value that would be important to detect

$$T_{\text{ES},1} = \inf\{t : \max_{t-m_1 \leq k \leq t-m_0} \max_{\mathcal{S} \subset [N]: |\mathcal{S}|=s} \sum_{(i,j) \in \mathcal{S}} U_{k,t,\delta}^{(i,j)} \geq b\},$$

- ▶ exist a recursive implementation (similar to CUSUM)
- ▶ for each possible \mathcal{S} , calculate

$$W_{\mathcal{S},t+1} = \max\{W_{\mathcal{S},t} + \sum_{(i,j) \in \mathcal{S}} U_{t,t+1,\delta}^{(i,j)}, 0\},$$

$$T_{\text{ES},1} = \inf\{t : \max_{\mathcal{S} \subset [N]: |\mathcal{S}|=s} W_{\mathcal{S},k} \geq b\}.$$

Exhaustive Search (ES) method (cont.)

- Approach 2: estimate p_1 for each hypothesize parameter values k and \mathcal{S}

$$\hat{p}_1(\mathcal{S}) = \frac{2}{|\mathcal{S}|(|\mathcal{S}| - 1)(t - k)} \sum_{(i,j) \in \mathcal{S}} \sum_{m=k+1}^t [X_m]_{ij},$$

$$T_{\text{ES},2} = \inf \left\{ t : \max_{t-m_1 \leq k \leq t-m_0} \max_{\mathcal{S} \subset [N]: |\mathcal{S}|=s} \sum_{(i,j) \in \mathcal{S}} U_{k,t,\hat{p}_1(\mathcal{S})}^{(i,j)} \geq b \right\}.$$

- no recursive implementation
- **limitation of ES:** \mathcal{S} unknown, have to search all possible subsets of $\{1, \dots, N\}$. Number of possible subsets $|\Omega| = 2^N$, **exponential** in N .

Mixture method

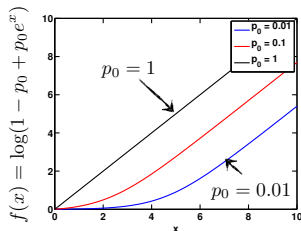
- ▶ exploit **structure**: typically community is a small subset
- ▶ assume two nodes (i, j) both in community with probability α
- ▶ α can be a guess for $|\mathcal{S}^*|/N$
- ▶ introduce indicator variable

$$Q_{ij} = \begin{cases} 1 & \text{w. p. } \alpha \\ 0 & \text{otherwise} \end{cases} \quad \forall i, j \in \mathcal{S}^*.$$

$$\ell(\kappa = k | p_1, \mathcal{S}) = \sum_{1 \leq i < j \leq N} \log \left\{ \mathbb{E}_{Q_{ij}} [(1 - Q_{ij}) + Q_{ij} \prod_{m=k+1}^t \frac{p_1^{[X_m]_{ij}} (1 - p_1)^{1 - [X_m]_{ij}}}{p_0^{[X_m]_{ij}} (1 - p_0)^{1 - [X_m]_{ij}}}] \right\} = \sum_{1 \leq i < j \leq N} h(U_{k,t,p_1}^{(i,j)}).$$

Mixture method (cont.)

$$h(x) \triangleq \log\{1 - \alpha + \alpha \exp(x)\}$$

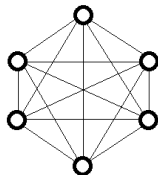


$$T_{\text{Mix}} = \inf\{t : \max_{t-m_1 \leq k \leq t-m_0} \sum_{1 \leq i < j \leq N} h(U_{k,t,\delta}^{(i,j)}) \geq b\},$$

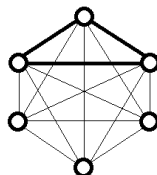
No search over subset \max_S .

Drawback of Mixture method

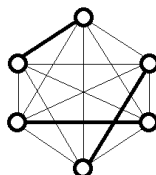
- ▶ statistics of Mixture method can be gathered from “false” community
- ▶ can increase false alarm rate



Social Network



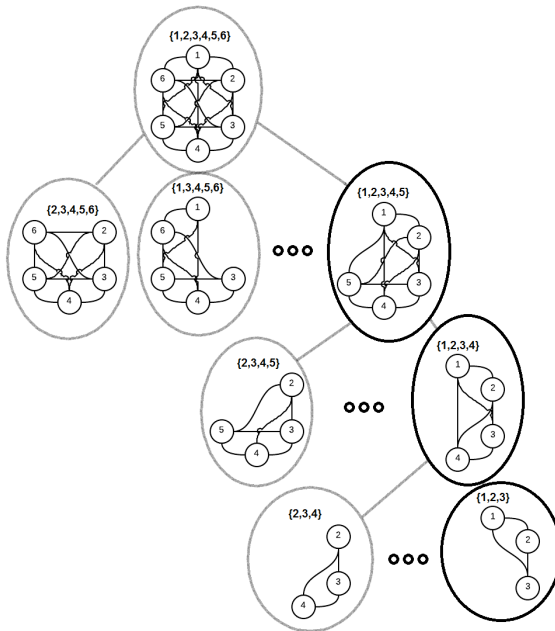
Social Network
with Subset



Mixture Model
Possible Outcome

- ▶ Hierarchical Mixture method (H-Mix) solves this problem by introducing dendrogram decomposition of the graph

Hierarchical Mixture method (H-Mix)



Algorithm 1 Hierarchical Mixture Method

- 1: Input: $\{X_m\}_{m=1}^t, X_m \in \mathbb{R}^{N \times N}$
 - 2: Output: $\{P_k\}_{k=1}^t \in \mathbb{R}^t$, a set of statistics for each hypothesized changepoint location k .
 - 3: **for** $k = 1 \rightarrow t$ **do**
 - 4: $\mathcal{S} = \llbracket N \rrbracket$
 - 5: **while** $|\mathcal{S}| > s$ **do**
 - 6: $i^* = \operatorname{argmax}_{i \in \mathcal{S}} M(\mathcal{S} \setminus \{i\})$
 - 7: $\mathcal{S} = \mathcal{S} \setminus \{i^*\}$
 - 8: **end while**
 - 9: $P_k = M(\mathcal{S})$
 - 10: **end for**
-

Complexity

Table : Complexities of algorithms under various conditions regarding k and N .

	$ \mathcal{S} \gg N/2$	$ \mathcal{S} \ll N/2$	$ \mathcal{S} \sim N/2$
Exhaustive Search	$\mathcal{O}(N^{N- \mathcal{S} })$	$\mathcal{O}(N^{ \mathcal{S} })$	$\mathcal{O}(2^{\frac{ \mathcal{S} }{2}})$
Mixture Model	$\mathcal{O}(N^2)$	$\mathcal{O}(N^2)$	$\mathcal{O}(N^2)$
Hierarchical Mixture	$\mathcal{O}(N^3)$	$\mathcal{O}(N^4)$	$\mathcal{O}(N^4)$

Choice of b

Choice of threshold b involves a tradeoff between **ARL** and **EDD**:

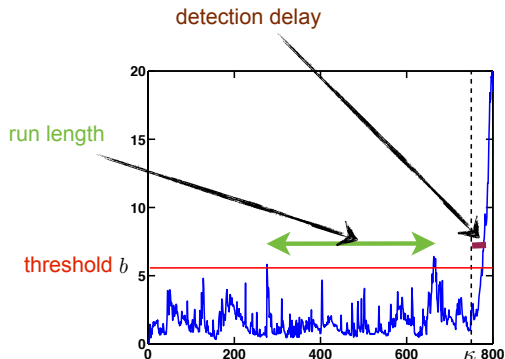
ARL (average run length) (captures false-alarm-rate)

- ▶ usually choose b to make ARL large $\sim 5000, 10000$
- ▶ for large N simulating ARL via Monte Carlo is hard
- ▶ accurate theoretical approximation for ARL is highly valuable

EDD (expected detection delay)

- ▶ a relatively small number ~ 10
- ▶ theoretical approximation provides useful insight

Performance metrics



- average run length (ARL):

$$\mathbb{E}^{\infty}\{T\}$$

- expected detection delay (EDD):

$$\sup_k \text{ess sup } \mathbb{E}^k\{T - k | T > k\}$$

Theoretical results

We obtain analytical expression for ARL of Mixture method

Theorem

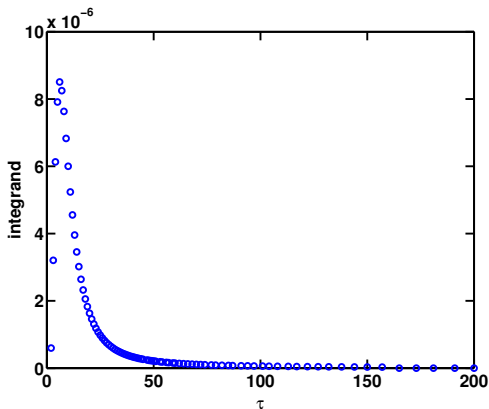
When $b \rightarrow \infty$, an *upper approximation* to the ARL $\mathbb{E}^\infty[T_{\text{mix}}]$ of the Mixture method with known p_1 is given by:

$$ARL_{\text{UA}} = \left[\int_{\sqrt{2N/m_1}}^{\sqrt{2N/m_0}} \frac{y\nu^2(y\sqrt{\gamma(\theta_y)})}{H(N, \theta_y)} dy \right]^{-1}, \quad (1)$$

and a *lower approximation* to the ARL is given by:

$$ARL_{\text{LA}} = \left[\sum_{\tau=m_0}^{m_1} \frac{2N\nu^2(2N\sqrt{\gamma(\theta_\tau)}/\tau^2)}{\tau^2 H(N, \theta_\tau)} \right]^{-1}, \quad (2)$$

- ▶ expressions can be evaluated explicitly
- ▶ no Monte Carlo simulation needed



only a few τ values play role in the summation

- ▶ accurate approximation
- ▶ can be used to determine threshold b for given ARL

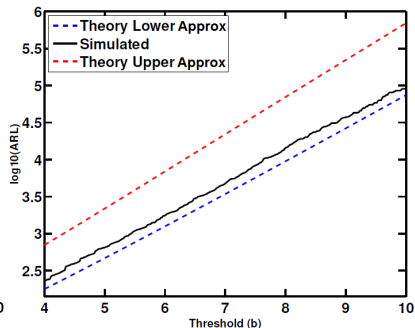
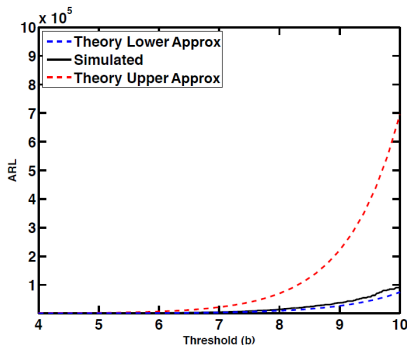
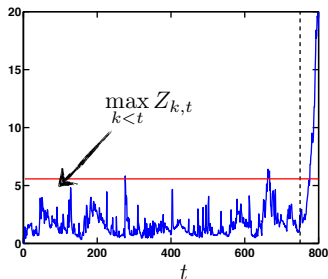


Table : Theoretical vs. simulated thresholds for $p_0 = 0.3$, $p_1 = 0.8$, and $N = 6$. The threshold b calculated using theory is very close to the corresponding threshold obtained using simulation.

ARL	Theory b	Simulated b
5000	7.37	7.04
10000	8.05	7.64

Proof techniques

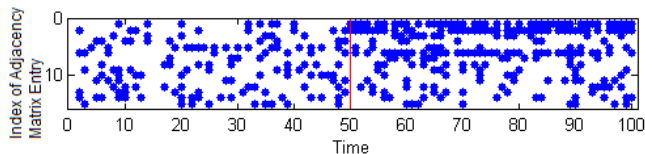
- ▶ detection statistic forms a random walk $\max_{k < t} Z_{k,t}$



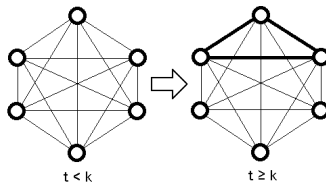
- ▶ to calculate ARL = calculate boundary hitting probability of a noise-like random walk
- ▶ the moment of detection of a stopping time, asymptotically exponentially distributed
- ▶ when $b \rightarrow \infty$, approximately $\mathbb{P}^\infty\{\max_{k < t} Z_{k,t} \geq b\} \approx m\lambda$

Numerical performance analysis

Detect emergence of a community



interact with $p_0 \Rightarrow$ interact with p_1



size of community is $|\mathcal{S}^*| = s$

Table : Comparison of detection delays for various cases when $N = 6$. The numbers inside the brackets are the threshold b such that $ARL = 5000$.

	$T_{ES,1}$ $\delta = p_1$	T_{Mix} $\delta = p_1$	T_{H-Mix}	T_{Mix} $\delta = p_1 - 0.1$
$s = 3, p_0 = 0.2, p_1 = 0.9$	3.8 (9.96)	4.3 (6.71)	3.8 (9.95)	6.0 (6.71)
$s = 3, p_0 = 0.3, p_1 = 0.7$	9.5 (10.17)	12.8 (6.77)	10.8 (10.18)	23.3 (6.77)
$s = 4, p_0 = 0.3, p_1 = 0.7$	5.0 (8.48)	6.7 (6.88)	6.4 (10.17)	11.0 (6.88)

Robustness against false communities

Table : ARL and DD for each algorithm under the conditions $p_0 = 0.2, p_1 = 0.9, k = 3$, and $N = 6$ where the $ARL = 5000$.

	Threshold	Detection Delay
$T_{ES,1}$	9.96	49.74
T_{Mix}	6.71	4.30
T_{H-Mix}	9.95	100.74

Mixture method incorrectly reacts to false community very quickly.

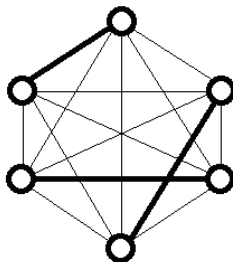
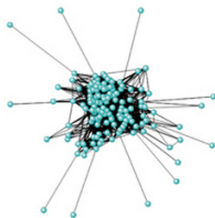


Table : Comparison of detection delays for a larger network $N = 50$. The numbers inside the brackets are the threshold b such that $\text{ARL} = 5000$.

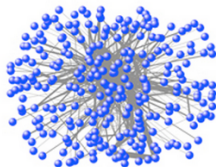
	$T_{\text{Mix}}, \delta = p_1$
$p_0 = 0.3, p_1 = 0.7, s = 10$	27.5 (-7.44)
$p_0 = 0.3, p_1 = 0.7, s = 20$	1.1 (-7.41)

Summary

- ▶ detect emergence of a community in sequential data
- ▶ present a new change-point detection approach
- ▶ three methods: exhaustive search (ES), mixture (Mix), and hierarchical mixture (H-Mix) methods, all able to detect the community quickly in different settings
 - ▶ complexity: $\text{Mix} < \text{H-Mix} \ll \text{ES}$
 - ▶ robustness: $\text{Mix} < \text{H-Mix} \approx \text{ES}$
- ▶ accurately theoretically characterize performance of Mix method
- ▶ future: apply on real data and larger networks: Enron data set



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