

## Signaling

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- An agent may be willing to reveal private information if he obtains greater utility
- Will not want to reveal information if the signal is too costly

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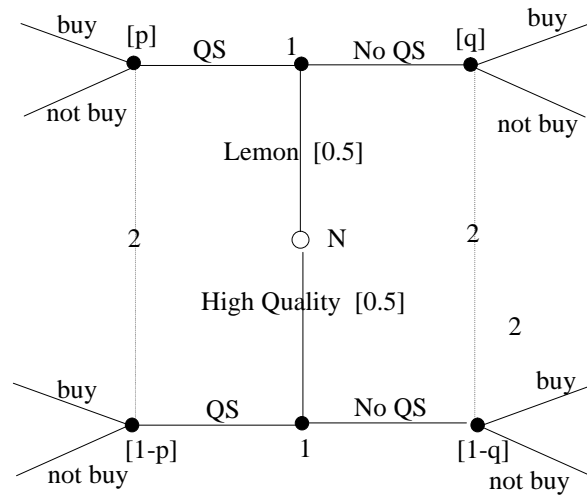
## Signaling Games

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- Nature draws a type  $t_i$  for the Sender from a set of feasible types  $T = \{t_1, \dots, t_n\}$  according to a probability distribution  $p(t_i)$ , where  $p(t_i) > 0$  for every  $i$  and  $p(t_1) + \dots + p(t_n) = 1$ .
- The Sender observes  $t_i$  and then chooses a message  $m_j$  from a set of feasible messages  $M = \{m_1, \dots, m_n\}$
- The Receiver observes  $m_j$  (but not  $t_i$ ) and then chooses an action  $a_k$  from a set of feasible actions  $A = \{a_1, \dots, a_n\}$ .
- Payoffs are given by  $U_s(t_i, m_j, a_k)$  and  $U_R(t_i, m_j, a_k)$

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## Signaling Example



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## Signaling Game Pure Strategies

- Sender:
  - Strategy 1: Play  $m_1$  if nature draws either type
  - Strategy 2: Play  $m_1$  if nature draws  $t_1$  and play  $m_2$  if nature draws  $t_2$
  - Strategy 3: Play  $m_2$  if nature draws  $t_1$  and play  $m_1$  if nature draws  $t_2$
  - Strategy 4: Play  $m_2$  if nature draws either type
- Receiver
  - Strategy 1: Play  $a_1$  no matter what message the Sender chooses
  - Strategy 2: Play  $a_1$  if the Sender chooses  $m_1$  and play  $a_2$  if the sender chooses  $m_2$
  - Strategy 3: Play  $a_2$  if the Sender chooses  $m_1$  and play  $a_1$  if the sender chooses  $m_2$
  - Strategy 4: Play  $a_2$  no matter what message the Sender chooses

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## Signaling Requirements

- Requirement 1: After observing any message  $m_j$  from  $M$ , the Receiver must have a belief about which types could have sent  $m_j$ . Denote this belief by the probability distribution  $\mu(t_i|m_j)$ , where  $\mu(t_i|m_j) \geq 0$  for each  $t_i$  in  $T$ , and  $\sum_{t_i \in T} \mu(t_i|m_j) = 1$ .

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## Signaling Requirements

- Requirement 2R: For each  $m_j$  in  $M$ , the Receiver's action  $a^*(m_j)$  must maximize the Receiver's expected utility, given the belief  $\mu(t_i|m_j)$  about which types could have sent  $m_j$ . That is,  $a^*(m_j)$  solves
$$\max_{a_k \in A} \sum_{t_i \in T} \mu(t_i|m_j) U_R(t_i, m_j, a_k).$$
- Requirement 2S: For each  $t_i$  in  $T$ , the Sender's message  $m^*(t_i)$  must maximize the Sender's utility, given the Receiver's strategy  $a^*(m_j)$ . That is,  $m^*(t_i)$  solves
$$\max_{m_j \in M} U_S(t_i, m_j, a^*(m_j)).$$

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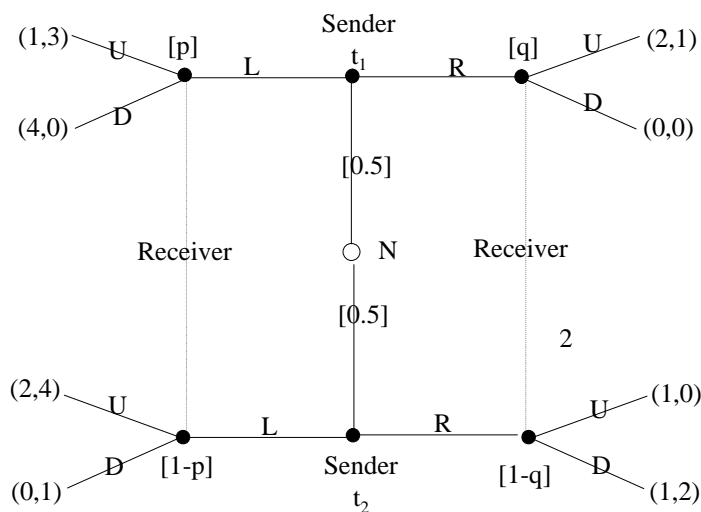
## Signaling Requirements

- Given the sender's strategy, let  $T_j$  denote the set of types that send the message  $m_j$
- Requirement 3: For each  $m_j$  in  $M$ , if there exists  $t_i$  in  $T$  such that  $m^*(t_i) = m_j$ , then the Receiver's belief at the information set corresponding to  $m_j$  must follow from Bayes' rule and the Sender's strategy:  

$$\mu(t_i | m_j) = p(t_i) / \sum_{t_i \in T_j} p(t_i).$$
- Definition: A pure-strategy perfect Bayesian equilibrium in a signaling game is a pair of strategies  $m^*(t_i)$  and  $a^*(m_j)$  and a belief  $\mu(t_i | m_j)$  satisfying Signaling Requirements (1), (2R), (2S), and (3).

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## Example



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## Possible Pure Strategy PBE

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- 1) Pooling on L
- 2) Pooling on R
- 3) Separation with  $t_1$  playing L and  $t_2$  player R
- 4) Separation with  $t_1$  playing R and  $t_2$  player L

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## Pooling on L

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- Suppose a Sender's strategy is (L,L)
  - Receiver's belief  $(p, 1-p)$  is on the equilibrium path, and her belief at this information set is determined by Bayes' rule and the Sender's strategy:  $p = .5$
  - Receiver's best response on the equilibrium path is to play U
  - Sender earns 1 (if  $t_1$ ) or 2 (if  $t_2$ )
- How would Receiver react to R?
  - If Receiver plays U to R, then Sender type  $t_1$  gets 2, which is better than he got by playing L.
  - If Receiver plays D to R, then Sender is better off by playing L instead
  - So if an pooling PBE (L,L) exists, Receiver must play (U,D) to L and R.

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## Pooling on L

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- When will a Receiver Play D to R?
  - She plays D if  $\Pi(D) \geq \Pi(U)$ , or  
 $q(0) + (1-q)(2) \geq q(1) + (1-q)(0) \rightarrow$   
 $q \leq 2/3$
- So  $[(L,L), (U,D), p=0.5, q]$  is a pooling PBE for any  $q \leq 2/3$

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## Pooling on R

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- Suppose a Sender's strategy is (R,R)
  - Then  $q = 0.5$
  - Receiver's best response to R is D
  - Sender earns 0 ( $t_1$ ) or 1 ( $t_2$ )
- But Sender type  $t_1$  could earn 1 by playing L (knowing Receiver's best response is U for any value of  $p$ )
- So there is no equilibrium in which the Sender plays (R,R)

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## Separation with $t_1$ playing L

- If Sender plays (L, R) then both of the Receiver's information sets are on the equilibrium path
  - So  $p = 1$  and  $q = 0$
  - Receiver's best response to these beliefs are U and D respectively
  - Both Sender types earn 1
- Given this Receiver response, is the Sender strategy optimal?
  - If type  $t_2$  deviates by playing L instead, then the Receiver responds with U, earning  $t_2$  a payoff of 2
- This separation equilibrium does not exist

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## Separation with $t_1$ playing R

- If Sender plays (R, L)
  - $p = 0$  and  $q = 1$
  - Receiver's best response is (U, U)
  - Both Sender types earn 2
- If  $t_1$  were to deviate by playing L
  - Receiver responds with U
  - $t_1$ 's payoff is 1
  - So no incentive to deviate
- If  $t_2$  were to deviate by playing R
  - Receiver responds with U
  - $t_2$ 's payoff is 1
  - So no incentive to deviate
- So  $[(R,L),(U,U),p=0,q=1]$  is a separating PBE

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