

Homework #5 - Solutions

Engineering Optimization (ISyE 4231) - Fall 2000

Problem 1.

Question 3. This schedule is feasible.

Aircraft 1 : $AT(1,4) = AT(1,3) = 1$ and $TT(4,3) = 1$

Aircraft 2 : $AT(2,8) = AT(2,2) = 1$ and $TT(8,2) = 1$

Aircraft 3 : $AT(3,1) = 1$

Aircraft 4 : $AT(4,7) = AT(4,6) = 1$ and $TT(7,6) = 1$.

Moreover, this schedule satisfies the maintenance restriction of aircraft 1 because $fly(4) + flight\ time\ from\ city\ 8\ to\ city\ 9 + fly(3) = 150 + 60 + 120 = 330 < 337$ and the ending time of aircraft 1's schedule = $298 + 150 = 448 < 630$.

Alternative feasible schedule : There are no alternative feasible schedules without subcontracting more aircraft. The reason is as follows.

Aircraft 3 : Aircraft 3 has to serve trip 1 since trip 1 is previously scheduled to aircraft 1.

No trips can be served immediately after or immediately before trip 1 by the same aircraft because $TT(2,1) = \dots = TT(8,1) = 0$ and $TT(1,2) = \dots = TT(1,8) = 0$.

So aircraft 3 can serve only trip 1 during this schedule period.

Aircraft 1 : From Question 4, we know that we need to subcontract one extra aircraft to serve trip 3, 5, 6 and 8. So we have to serve trip 4 and trip 7 with our aircraft not to subcontract more aircraft. Since $AT(1,4) = 1$ and $AT(2,4) = AT(3,4) = AT(4,4) = 0$, Aircraft 1 has to serve trip 4. Because of the maintenance restriction of aircraft 1, trip 3 is the only possible next trip that can be served by aircraft 1. Aircraft 1 has to serve trip 3. (Suppose not. We need to subcontract more aircraft because only aircraft 2 and 4 are available for serving trip 3, 5, 6 and 8.) So aircraft 1 has to serve trip 4 and trip 3.

Aircraft 4 : Aircraft 4 has to serve trip 7. (Suppose not we need to subcontract one more aircraft for trip 7 because $AT(4,7) = 1$ and $AT(1,7) = AT(2,7) = AT(3,7) = 0$). Only trip 6 can be served immediately after trip 7 out of trip 5, 6, 8 from the TT matrix. So aircraft 4 has to serve trip 6. (Suppose not we need to subcontract more aircraft because only aircraft 2 is available for trip 5, 6 and 8.) So aircraft 4 has to serve trip 7 and trip 6.

Aircraft 2 : Aircraft 2 has to serve trip 2 since trip 2 is previously scheduled to aircraft 2. Aircraft 2 has to serve trip 8. (Suppose not No aircraft are available for trip 5 and 8 because $TT(5,2) = 0$.) So aircraft 2 has to serve trip 8 and trip 2.

Therefore we have to subcontract one extra aircraft for serving trip 5.

Obviously, there are several alternative feasible solutions with subcontracting more aircraft.

Question 4. No. Consider trips 1, 3, 5 and 8. Since these trips all intersect each other (see figure), we need 4 different aircraft to serve them. Furthermore, trip 6 intersects trips 1, 3, and 5, so the aircraft serving trips 1, 3, and 5 cannot serve trip 6. Finally, there is not enough time to fly from the destination of trip 8 to the departure location of trip 6 on time, so if an aircraft serves trip 8 then it cannot serve trip 6. Therefore, among the trips 1, 3, 5, 6 and 8, at least one should be subcontracted.

More formally, we have

$$TT(1,3) = TT(1,5) = TT(1,6) = TT(1,8) = 0$$

$$TT(3,1) = TT(3,5) = TT(3,6) = TT(3,8) = 0$$

$$TT(5,1) = TT(5,3) = TT(5,6) = TT(5,8) = 0$$

$$TT(6,1) = TT(6,3) = TT(6,5) = TT(6,8) = 0$$

$$TT(8,1) = TT(8,3) = TT(8,5) = TT(8,6) = 0.$$

This means that we need different airplanes to serve each trip 1, trip 3, trip 5, trip 6 and trip 8. So there are no feasible schedules without subcontracting extra aircraft (we have 4 airplanes and any two of above five trips cannot be served by one aircraft.)

Problem 2.

(a) Standard form

$$\begin{aligned} \max \quad & 3x_1 + x_2 \\ \text{s.t.} \quad & -2x_1 + x_2 + x_3 = 2 \\ & x_1 + x_2 + x_4 = 6 \\ & x_1 + x_5 = 4 \\ & x_1, x_2, x_3, x_4, x_5 \geq 0 \end{aligned}$$

(b) Above standard-form linear program has the following coefficient data.

	x_1	x_2	x_3	x_4	x_5	
Max c	3	1	0	0	0	b
	-2	1	1	0	0	2
A	1	1	0	1	0	6
	1	0	0	0	1	4

$x^{(0)} = (0, 0, 2, 6, 4)$ and x_3, x_4 and x_5 are basic.

First iteration

Simplex directions ($A\Delta x = 0$)

Simplex direction increasing x_1 when Δx_1 is fixed =1.:

$$\begin{aligned} -2(1) + 1(0) + 1\Delta x_3 + 0\Delta x_4 + 0\Delta x_5 &= 0 \\ +1(1) + 1(0) + 0\Delta x_3 + 1\Delta x_4 + 0\Delta x_5 &= 0 \\ +1(1) + 0(0) + 0\Delta x_3 + 0\Delta x_4 + 1\Delta x_5 &= 0 \end{aligned}$$

By solving above equation system we have $\Delta x = (1, 0, 2, -1, -1)$

Simplex direction increasing x_2 when Δx_2 is fixed =1.:

$$\begin{aligned} -2(0) + 1(1) + 1\Delta x_3 + 0\Delta x_4 + 0\Delta x_5 &= 0 \\ +1(0) + 1(1) + 0\Delta x_3 + 1\Delta x_4 + 0\Delta x_5 &= 0 \\ +1(0) + 0(1) + 0\Delta x_3 + 0\Delta x_4 + 1\Delta x_5 &= 0 \end{aligned}$$

By solving above equation system we have $\Delta x = (0, 1, -1, -1, 0)$

Improving simplex directions

$\bar{c}_1 = (3, 1, 0, 0, 0) \bullet (1, 0, 2, -1, -1) = 3 > 0$, $\Delta x = (1, 0, 2, -1, -1)$ is a improving simplex direction

$\bar{c}_2 = (3, 1, 0, 0, 0) \bullet (0, 1, -1, -1, 0) = 1 > 0$, $\Delta x = (0, 1, -1, -1, 0)$ is a improving simplex direction

Entering variable : candidate entering variables are x_1 and x_2 . Let's pick x_1 .

Leaving variable : $\min\{6/1, 4/1\} = 4$. minimized for x_5 . x_5 leaves the basis.

New basic feasible solution :

$$\begin{aligned} x^{(1)} &\leftarrow x^{(0)} + \mathbf{I}\Delta x = (0, 0, 2, 6, 4) + 4(1, 0, 2, -1, -1) = (4, 0, 10, 2, 0) \\ &\text{and } x_1, x_3 \text{ and } x_4 \text{ are basic} \end{aligned}$$

Second iteration

Simplex directions ($A\Delta x = 0$)

Simplex direction increasing x_2 when Δx_2 is fixed =1.:

$$-2\Delta x_1 + 1(1) + 1\Delta x_3 + 0\Delta x_4 + 0(0) = 0$$

$$+ 1\Delta x_1 + 1(1) + 0\Delta x_3 + 1\Delta x_4 + 0(0) = 0$$

$$+ 1\Delta x_1 + 0(1) + 0\Delta x_3 + 0\Delta x_4 + 1(0) = 0$$

By solving above equation system we have $\Delta x = (0, 1, -1, -1, 0)$

Simplex direction increasing x_5 when Δx_5 is fixed =1.:

$$-2\Delta x_1 + 1(0) + 1\Delta x_3 + 0\Delta x_4 + 0(1) = 0$$

$$+ 1\Delta x_1 + 1(0) + 0\Delta x_3 + 1\Delta x_4 + 0(1) = 0$$

$$+ 1\Delta x_1 + 0(0) + 0\Delta x_3 + 0\Delta x_4 + 1(1) = 0$$

By solving above equation system we have $\Delta x = (-1, 0, -2, 1, 1)$

Improving simplex directions

$\bar{c}_2 = (3, 1, 0, 0, 0) \bullet (0, 1, -1, -1, 0) = 1 > 0$, $\Delta x = (0, 1, -1, -1, 0)$ is a improving simplex direction

$\bar{c}_5 = (3, 1, 0, 0, 0) \bullet (-1, 0, -2, 1, 1) = -3 < 0$, $\Delta x = (-1, 0, -2, 1, 1)$ is **not** a improving simplex direction

Entering variable : x_2 .

Leaving variable : $\min\{10/1, 2/1\} = 2$. minimized for x_4 . x_4 leaves the basis.

New basic feasible solution :

$$x^{(2)} \leftarrow x^{(1)} + \mathbf{I}\Delta x = (4, 0, 10, 2, 0) + 2(0, 1, -1, -1, 0) = (4, 2, 8, 0, 0)$$

and x_1, x_2 and x_3 are basic.

Third iteration

Simplex directions ($A\Delta x = 0$)

Simplex direction increasing x_4 when Δx_4 is fixed =1.:

$$-2\Delta x_1 + 1\Delta x_2 + 1\Delta x_3 + 0(1) + 0(0) = 0$$

$$+ 1\Delta x_1 + 1\Delta x_2 + 0\Delta x_3 + 1(1) + 0(0) = 0$$

$$+ 1\Delta x_1 + 0\Delta x_2 + 0\Delta x_3 + 0(1) + 1(0) = 0$$

By solving above equation system we have $\Delta x = (0, -1, 1, 1, 0)$

Simplex direction increasing x_5 when Δx_5 is fixed =1.:

$$-2\Delta x_1 + 1\Delta x_2 + 1\Delta x_3 + 0(0) + 0(1) = 0$$

$$+ 1\Delta x_1 + 1\Delta x_2 + 0\Delta x_3 + 1(0) + 0(1) = 0$$

$$+ 1\Delta x_1 + 0\Delta x_2 + 0\Delta x_3 + 0(0) + 1(1) = 0$$

By solving above equation system we have $\Delta x = (-1, 1, -3, 0, 1)$

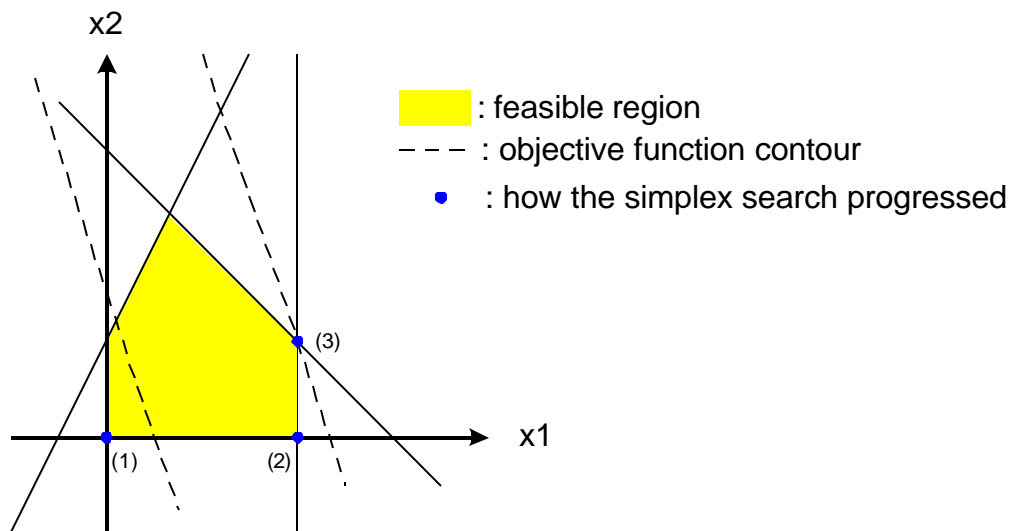
Improving simplex directions

$\bar{c}_2 = (3, 1, 0, 0, 0) \bullet (0, -1, 1, 1, 0) = -1 < 0$, $\Delta x = (0, -1, 1, 1, 0)$ is **not** a improving simplex direction

$\bar{c}_5 = (3, 1, 0, 0, 0) \bullet (-1, 1, -3, 0, 1) = -2 < 0$, $\Delta x = (-1, 1, -3, 0, 1)$ is **not** a improving simplex direction

There are no improving simplex directions. So the current solution is optimal.

(c)



Problem 3.

- (a) Yes
- (b) No. (E is not a extreme point.)
- (c) No. (H is not a extreme point.)
- (d) No. (Objective function values of sequence of solutions followed by simplex algorithm is non-decreasing in maximization problem and non-increasing in minimization problem. However, objective function values of this sequence is not monotone.)
- (e) No. (From A, F can not be visited without visiting G in simplex algorithm)